

A PROBABILISTIC ANALYSIS OF THE DYNAMICS OF STOCHASTIC EQUATIONS DETERMINED BY THE GILLESPIE ALGORITHM

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1. INTRODUCTION

The accurate modelling of the time dynamics of discrete physical systems is integral across the sciences. Such models often involve an approximation of both the system state space as a continuous quantity and the evolution of the system as a deterministic quantity. In systems with a small number of elements, such a model has limited validity, due to the increasing effect of stochasticity in the dynamics, and the increasing granularity of the state space. The Gillespie Algorithm provides a way of modelling these systems in such a way that neither of the aforementioned approximations are necessary. However, this implies a loss of determinism. In this essay, we consider to what extent we can forecast the evolution of Gillespie Algorithm derived models of such systems, and determine behaviours that are independent of inter-trial stochasticity.

Biochemical processes, and, in particular, gene regulatory networks, form a critical class of systems that are amenable to stochastic modelling through the Gillespie Algorithm. Such systems evolve through discrete reaction events, in which both the choice of reaction and the time between reactions are random quantities. In gene regulatory network models, there is no clear concept of mass conservation. In other words, the gain of one system substrate does not necessarily imply the loss of others, and vice versa. This reflects the fact the substrates of such systems are generally messenger RNA molecules, or proteins. The rate of change of these factors may be a product of the state of the system, but the actual production and degradation processes can be carried out by biological machinery external to the system. The Circadian Rhythm Model described in [GOL03] exhibits this property, as does our simplification of the model presented in Appendix B. For this reason, we concentrate on the class of models in which each discrete system evolution takes the form of a gain or loss of one substrate element.

The vast majority of the mathematical machinery we construct is unique to this essay, and is introduced as it comes up, so as not to detract from the overall flow of the exposition. In order that the reader does not get lost in a morass of definitions, we provide an index of terminology, in Appendix D.

Shortly prior to submission of this essay, it was found that our definition of a time-stable Gillespie Simulation is equivalent to that of an ergodic Markov Chain, and that many of the results presented in Chapter 4 and onwards, have generalised counterparts in [FAY95]. However, the proofs of these results are in general unique to this exposition, having been arrived at independently, and it is indicated if this is not the case. Reconstruction of the mathematical apparatus and nomenclature to coincide with that presented in [FAY95], and general Markov Chain theory, was not considered a worthwhile endeavour.

2. THE GILLESPIE ALGORITHM

In this section, we construct some of the basic mathematical objects necessary for our subsequent analysis, and provide the explicit form of the Gillespie Algorithm

Definition 2.1. Substrate State Space: We define the substrate state space, X , of a physical system with n distinct substrate types, as the set of allowable substrate vectors. So an element $x \in X$ is an n -dimensional vector, with the value of x_i denoting the quantity of the i^{th} substrate of the system. We define the **mass** of x to denote the sum of the numbers of its substrate elements. In this essay, we consider physical systems where substrate levels take positive integer values. Therefore, $X \cong \mathbb{N}^n$.¹ In particular we define, the i^{th} **basis** vector of X : e^i , as follows:

$$e_j^i = \delta_{ij}$$

Here, δ signifies the Kronecker delta function.

Definition 2.2. System State Space We define the system state space Y as the product space $X \times \mathbb{R}^+$. A system state $y \in Y$ is of the form (x, t) , with $x \in X$, $t \in \mathbb{R}^+$, and describes the system having substrate levels corresponding to x , at time t .

Definition 2.3. Reaction State Space: We consider physical systems that evolve through the occurrence of m separate reactions, $r_1 \dots r_m$. We define the reaction state space Q of the system as the discrete set consisting of these reactions. Each reaction can be considered as a function mapping X to itself. This essay is chiefly concerned with systems where the following assumptions apply:

- (1) $m = 2n$
- (2) $r_i(x) = x + e^i \quad i \in \{1, 2, \dots, n\}$
- (3) $r_{i+n}(x) = x - e^i \quad i \in \{1, 2, \dots, n\}$

Definition 2.4. Reaction Rate Vector Field: The reaction rate vector field, R , of a physical system, is a function

$$R : Y \mapsto \mathbb{R}^m$$

The value $R_i(y)$ denotes the rate of occurrence of reaction r_i in the system, given system state y . Note that reaction rates can only take non-negative values, ie $R_i(y) \geq 0, \forall y \in Y$. In this essay, we will be dealing solely with physical systems whose reaction rate vector fields are independent of time, and therefore purely a function of the current substrate levels. Therefore, we will henceforth consider R as a function with effective domain X . In other words, given $y \in Y$ such that $y = (x, t)$ we take $R(y) = R(x)$, and use the two terms interchangeably. In this essay, the following assumption applies:

- (1) $R_i(x) \neq 0 \quad \forall i \in \{1, \dots, n\}$
- (2) $R_i(x) = 0$ iff $x_i = 0, \forall i \in \{n + 1, \dots, 2n\}$

¹Here, we consider \mathbb{N} to include 0

Definition 2.5. Gillespie Model: Define a Gillespie model as the pair (X, R) consisting of a substrate state space, and reaction rate vector field. Note that any two physical systems sharing the same Gillespie model have identical dynamics. In this sense, we can consider a Gillespie model to uniquely define a physical system being modelled.

We are now in a position to define the Gillespie Algorithm, which generates a stochastic trajectory of the system state of a Gillespie model in time. In the exposition of the algorithm to come, the following notation is used:

Notation 2.6.

$$(1) \lambda(x) = \sum_{i=1}^{i=m} R_i(x) = \|R(x)\|_{L_1} \quad (\text{as every rate is by definition non-negative}).$$

λ corresponds to the total reaction rate of the system.

(2) I is the iteration counter. On the i^{th} iteration of the Gillespie Algorithm, $I := i$

Algorithm 2.7. Classical Gillespie Algorithm [WIL06]

(1) **Initialisation:** Set an initial system state $y_0 = (x_{t_0}, t_0)$, and calculate $R(x_{t_0})$. Set $I = 1$

(2) **Random Number Generation and Scaling:** Calculate $\lambda(x_{t_{I-1}})$. Simulate a realisation of the following random variables

$$u_I \sim U[0, 1] \quad w_I \sim \exp(1)$$

Define:

$$\tilde{w}_I = \frac{1}{\lambda(x_{t_{I-1}})} \times w_i \quad \tilde{u}_I = u_I \times \lambda(x_{t_{I-1}})$$

Note that $\tilde{w}_I \sim \exp(\lambda(x_{t_{I-1}}))$, and $\tilde{u}_I \sim U[0, \lambda(x_{t_{I-1}}))$

(3) **System Update:** We define the set of intervals $\{c_k\}_{i=1}^m$

$$c_k = \sum_{i=1}^{k-1} R_i(x_{t_{I-1}}) + [0, R_k(x_{t_{I-1}}))$$

. Note that:

$$(a) \quad \bigcup_{i=1}^m c_i = [0, \lambda(x_{t_{I-1}}))$$

$$(b) \quad c_i \cap c_j = \emptyset \quad \forall i, j$$

Therefore \tilde{u}_I is in exactly one of the set of intervals $\{c_k\}$. Call this the j^{th} interval. Set

$$y_I = \left(r_j(x_{t_{I-1}}), t_I + w_I \right)$$

(4) **Reset:** Store y_I . Set $I = I + 1$. Return to step (2)

Explanation

The algorithm models each reaction as a poisson process, of rate equal to its respective reaction rate. The sum of reactions is therefore a poisson process with rate equal to the sum of the rates of the respective reactions. Waiting times between reactions are therefore exponentially distributed, and given by the random variables of the form $\{\tilde{w}_I\}$. After each waiting time, we choose a reaction to simulate. The probability of choosing a reaction is equal to its reaction rate as a proportion of the total reaction rate of the system.

Definition 2.8. Define a ***K* step Gillespie Simulation** on a Gillespie Model as the sequence: $\{y_i\}_{i=0}^{i=K}$, obtained by performing the Gillespie algorithm K times given an initial system state y_0 . A Gillespie simulation with no step number specified is assumed to be an infinite sequence. Note that simulations are not well defined, due to the inherent stochasticity of the algorithm.

Definition 2.9. Define the **trajectory**, $x(t) : (t_0, T) \rightarrow X$, of a Gillespie Simulation $\{y_i\}_{i=0}^{\infty}$ as follows:

$$x(t) = x_{t_i} \quad t \in [t_i, t_{i+1})$$

Define the **initial condition** of a Gillespie trajectory as the state $y_0 = (x_0, t_0)$. Here, $T = \lim_{i \rightarrow \infty} t_i$ for an infinite Gillespie simulation, and $T = t_K$ in the obvious (omitted) modification of the definition to accommodate a K -step Gillespie Simulation.

In essence, when we iteratively perform the Gillespie algorithm on a Gillespie Model with initial conditions, we get a discrete set of vectors $\{y_i\}$, which capture every single change occurring in the system during application of the algorithm. This is the Gillespie simulation. The trajectory continuously models the actual time evolution of the system, which is unchanging except at the set of timepoints recorded by the Gillespie simulation.

Definition 2.10. Given a point $x(t')$ on a trajectory, we name t' as the **elapsed (external) time**. We also define the **internal time** p of a point $x(t')$ on the trajectory as follows:

$$p[x(t')] = i : t' \in [t_i, t_{i+1})$$

Note that the space of internal times is isomorphic to \mathbb{N} . The internal time corresponds to

the number of changes the system state has gone through. We denote the space of internal times as \mathbb{T}

Definition 2.11. Random Sequence Space:

Realisation of an infinitely long Gillespie simulation involves generation of the following sequences of random variables:

$$\{u_i\}_{i=1}^{\infty} : u_j \sim U[0, 1] \quad \forall j \in \mathbb{N}$$

$$\{w_i\}_{i=1}^{\infty} : w_j \sim \text{exp}(1) \quad \forall j \in \mathbb{N}$$

We accordingly define the random sequence space Ω as follows:

$$\Omega = U \times W$$

$$U = \{u : u = \{u_i\}_{i=1}^{\infty}\}; \quad u_i \sim U[0, 1]$$

$$W = \{w : w = \{w_i\}_{i=1}^{\infty}\}; \quad w_i \sim \text{exp}(1)$$

Proposition 2.12. A Gillespie Simulation (and thereby a trajectory) on a Gillespie Model can be uniquely determined, given the pair: (x_0, ω) , where $x_0 \in X$, and $\omega \in \Omega$.

Proof. Decompose ω into the pair of sequences (u, w) . The proposition is equivalent to stating that a Gillespie Simulation is deterministic, given (x_0, u, w) . All of the randomness in the Gillespie algorithm is associated with the generation of the sequences: w and u . Therefore, the result is trivial. We have purposely excluded the provision of an initial elapsed timepoint t_0 in the initial condition (x_0, u, w) , since it has no effect on the dynamics of the Gillespie simulation: and merely serves to translate the simulation in time. \square

For clarity, we outline an explicit, modified version of the algorithm, which runs deterministically, given a triple (x_0, w, u) , below.

Algorithm 2.13. Modified Gillespie Algorithm

- (1) **Initialisation:** Generate an element $\omega \in \Omega$
Set initial system state and rates vectors: y_0 and $R(y_0)$, and iteration counter $I = 1$
- (2) **Random Number Scaling:** As in Algorithm 2.7
- (3) **System Update:** As in Algorithm 2.7
- (4) **Return to Step 2**

3. THE FORMULATION OF GILLESPIE SIMULATIONS AS RANDOM DYNAMICAL SYSTEMS

Definition 3.1. Random Dynamical System [LUD98] *A measurable random dynamical system on the measurable space (X, \mathcal{B}) over a metric dynamical system $(\Omega, \mathcal{F}, \mathbb{P}, (\theta(t))_{t \in \mathbb{T}})$ with time \mathbb{T} is a mapping:*

$$\varphi : \mathbb{T} \times \Omega \times X \rightarrow X, \quad (t, \omega, x) \mapsto \varphi(t, \omega, x)$$

with the following properties:

- (1) **Measurability:** φ is $\mathcal{B}(\mathbb{T}) \otimes \mathbb{F} \otimes \mathcal{B}$, \mathcal{B} -measurable
- (2) **Cocycle Property:** The mappings $\varphi(t, \omega) := \varphi(t, \omega, \cdot) : X \rightarrow X$ form a cocycle over $\theta(\cdot)$ ie they satisfy:

$$\varphi(0, \omega) = id_X \quad \forall \omega \in \Omega \tag{3.1}$$

$$\varphi(t + s, \omega) = \varphi(t, \theta(s)\omega) \circ \varphi(s, \omega) \quad \forall s, t \in \mathbb{T} \quad \omega \in \Omega \tag{3.2}$$

Here, we use ‘ \circ ’ to denote composition, which canonically defines an action on the left of the semigroup of self mappings, $X \rightarrow X$.

Theorem 3.2. A Gillespie Simulation on a Gillespie Model can be formulated as a Random Dynamical System

The proof of this theorem forms the crux of this chapter. We commence by defining the relevant spaces and measured espoused in Definition 3.1. In this chapter, we take for granted the assumptions of Definitions 2.3 and 2.4.

Definition 3.3. *We define a σ -algebra, \mathcal{B} , on the system state space Y . Recall that $Y \cong \mathbb{N}^n \times \mathbb{R}^+$ (Definition 2.2). We take \mathcal{B} to be the product of the Counting and Lebesgue measures on these respective spaces. \mathcal{B} is well defined, due to existence and uniqueness of the product measure [RUD87].*

Definition 3.4. *We define $(\Omega, \mathcal{F}, \mathcal{V})$ to be the measurable space such that:*

- (1) Ω is defined as in 2.11
- (2) \mathcal{F} is the σ -algebra generated by sets of the form $\{(u, w) : (u_i, w_i) \in I_i \times J_i \subseteq [0, 1] \times \mathbb{R}^+\}$, where I_i and J_i are open intervals, for some $i \in \mathbb{N}$. Since this is a countable product of Borel σ -algebras (which are generated by open intervals) on the support of each sequence element, we rest assured that it too is a σ -algebra ([RUD87]).

$$(3) \mathcal{V}(A) = \prod_{i=1}^{\infty} P((u_i, w_i) \in A_i).$$

Here the probability measure, P , is that derived from the probability density functions of u_i and w_i . \mathcal{V} , as a countable product of probability measures, is itself a probability measure [RUD87]

Definition 3.5. For a given Gillespie Model, we define the ‘flow’ map $\varphi : \mathbb{T} \times \Omega \times Y \mapsto Y$, as follows:

$$\varphi(d, (u, w), y_0) = (y_d)$$

Here, y_d is the system state obtained by conducting a d -step Gillespie Simulation on the triple (y_0, u, w) , which we have previously shown to be deterministic (2.12). $d \in \mathbb{T}$, the space of internal times, which we have previously noted is isomorphic to \mathbb{N} .

Note that 2.13 provides an explicit method of computation for the flow map.

Definition 3.6. We define a Borel σ -algebra on $\mathbb{T} : \mathcal{B}(\mathbb{T})$, by giving \mathbb{T} the discrete topology.

Note that every subset of \mathbb{T} is measurable, due to countability of the space combined with the fact that every discrete element of the space is measurable.

Theorem 3.7. *The Flow map, φ , is measurable*

Our plan of action is to decompose the flow map φ into the composition of two functions. One transforms the sequence u of uniform random variables into a sequence of reaction types, and the next converts the initial substrate vector to the final one using this reaction sequence, while calculating the reaction timepoint. The measurability of both of these functions then implies measurability of their composition, φ . Note that this decomposition is employed for purely theoretical purposes. Applying it algorithmically would be computationally expensive

Definition 3.8. Consider a Gillespie simulation on a model with m reaction types, denoted $r_1 \dots r_m$. We define S as the space of reaction sequences.

$$s \in S \Rightarrow s = \{s_i\}_{i=1}^{\infty}; \quad s_i \in Q, \quad \forall i$$

We define a measure on S in several stages: We first formalise the notion of distance between two reaction sequences. Let r_i, r_j be two reactions. Then define

$$|r_i - r_j| = \delta_{ij}$$

Here, δ_{ij} denotes the Kronecker Delta function. Now for arbitrary $q, s \in S$:

$$d(q, s) = \sum_{i=1}^{\infty} \frac{|q_i - s_i|}{2^i}$$

Such metrics on a sequence space are common in fields such as Dynamical Systems, so we neglect to prove that d does indeed define a metric, and refer the reader to, for example [DEV92]. Now that we have a metric, and therefore an induced metrizable topology, on S , we note the existence of a Borel σ -algebra $\mathcal{B}(S)$ generated by open sets in the topology. In other words, $\mathcal{B}(S)$ is generated by balls B of arbitrary radius in S . Note the following:

$$B\left(s, 2^{-k}\right) = \left\{q \in S : \{q_i\}_{i=1}^k = \{s_i\}_{i=1}^k\right\}$$

Definition 3.9. We define the ‘reaction sequence generating’ function of a Gillespie Model as follows:

$$\begin{aligned} \psi : \mathbb{T} \times \Omega \times Y &\mapsto (\mathbb{T} \times W \times S \times Y), \\ \psi(d, (u, w), y) &= (d, w, s, y) \end{aligned}$$

Here s is the sequence of reaction types occurring in a Gillespie simulation on the model given initial conditions (x, w, u) . Applying Algorithm 2.13 to this triple and storing only the generated reaction sequence would explicitly compute the function. Note that this is essentially a function mapping U to S : it acts as the identity function on other subspaces of its domain, and they are included purely in order that it can transmit information to functions it is composed with.

Definition 3.10. For a given Gillespie simulation, we define the function ϕ as follows

$$\begin{aligned} \phi : \mathbb{T} \times W \times S \times Y &\mapsto Y, \\ \phi(d, w, y_0) &= y_d \end{aligned}$$

Here, y_d is the system state reached by performing a d -step Gillespie simulation on y_0 , with the reaction sequence of the simulation specified by s , and the reaction timepoints by w .

The proofs of results leading to Theorem 3.7 are simplified by the following observations:

- (1) $\varphi(d, \omega, y) = \phi \circ \psi(d, \omega, y)$, so measurability of ϕ and ψ is a sufficient condition for measurability of φ
- (2) To prove measurability of a function $f : X \rightarrow Y$, with (Y, \mathcal{B}) a measurable space, it is sufficient to prove measurability of arbitrary elements of a generating set for the σ -algebra \mathcal{B} . Here we define a generating set of \mathcal{B} to be a set that generates the whole of \mathcal{B} under σ -algebra operations. [RUD87]

- (3) The σ -algebra that allows implementation of the Lebesgue measure on \mathbb{R}^+ is generated by open intervals. The σ -algebra required for implementation of the counting measure on a countable space is generated by discrete elements of that space. [RUD87]
- (4) The σ -algebra \mathcal{B} on Y is generated by the product of single elements of X , and open intervals on \mathbb{R}^+ . To prove measurability of a function f that maps onto Y , we therefore need only prove measurability of $f^{-1}(x, E)$, for an arbitrary element $x \in X$, and an arbitrary interval $E \in \mathbb{R}^+$

Lemma 3.11. *ψ is $\mathcal{B}(\mathbb{T}) \otimes \mathcal{F} \otimes \mathcal{B}$ measurable*

Not that this is equivalent to showing that $\psi^{-1}(d, \omega, A, y)$ is $\mathcal{B}(\mathbb{T}) \otimes \mathcal{F} \otimes \mathcal{B}$ measurable, for any discrete points $d \in \mathbb{T}, \omega \in \Omega, y \in Y$, and open interval $A \in S$.

Proof. The restriction of the range of ψ to $\mathbb{T} \times \Omega \times Y$ is the identity function, so it is sufficient to prove that $\psi^{-1}(A, \cdot)$ is measurable. Decomposing the product spaces Y and Ω into their constituent factor spaces, we have:

$$\begin{aligned} \psi^{-1}(A, \cdot) &= \{x \in X, t \in \mathbb{R}^+, d \in \mathbb{T}, u \in U, w \in W \text{ such that } \psi|_S(d, u, w, x, t) = A\} \\ &= \mathbb{R}^+ \times W \times \bigcup_{x \in X} \bigcup_{d \in \mathbb{T}} \left\{ u : \psi|_S(d, u, x, \cdot) = A \right\} \end{aligned}$$

The second equality holds due to the fact that internal and elapsed time calculations are independent of reaction sequence choice in the Gillespie Algorithm. Countability of X and \mathbb{T} imply that the above double union is countable. It remains to show measurability, for fixed $d \in \mathbb{T}, x \in X$, of:

$$\{u \in U : \psi|_S(d, u, x, \cdot) = A\}$$

Any open set in S is a countable union of open balls, so we may assume without loss of generality that A is an open ball. From 3.8, we see that

$$A = \{q \in S : \{q_i\}_{i=1}^k = \{s_i\}_{i=1}^k\}$$

for some fixed sequence of reactions $\{s_i\}_{i=1}^k$.

We know that $s_1 = r_j$, for some $j \in \{1, \dots, m\}$. For a Gillespie simulation with arbitrary initial condition y_0 , the first reaction is r_j if and only if u_1 falls in some interval, of which the size, and start/end points are determined by the computations in 2.13, and the current state of the substrate vector. Let us denote this interval I_{s_1} . If the j^{th} reaction rate is 0, then $I_{s_1} = \emptyset$. We inductively define I_{s_i} similarly. Then

$$\{u \in U : \psi|_S(u, x, d, \cdot) \in A\} = \prod_{i=1}^k \{u : u_i \in I_{s_i}\}$$

From Definition 3.4, we see that this is measurable, since, although we have not analysed the dimensions of each individual I_{s_i} , we know each one represents either an interval or the empty set. \square

Lemma 3.12. *ϕ is measurable*

Note that this is equivalent to showing that $\phi^{-1}(x, E)$ is $\mathcal{B}(\mathbb{T}) \otimes \mathcal{F} \otimes \mathcal{B}(\mathbb{S}) \otimes \mathcal{B}$ measurable for any $x \in X$ and interval $E \subseteq \mathbb{R}^+$, as (Y, \mathcal{B}) is generated by sets of this form.

Proof.

$$\phi^{-1}(x, E) = \bigcup_{x' \in X} \bigcup_{d \in \mathbb{T}} \{w \in W, t \in \mathbb{R}^+, s \in S : \phi(d, (x', t), w, s) \in (x, E)\}$$

Each of these unions is countable due to the countability of X and \mathbb{T} . Therefore, measurability of the whole set is equivalent to that of each component element of the double union.

Note that we may assume $t = 0$, without loss of generality, in the above equality. Changing the initial elapsed time value merely translates the final elapsed time value, so does not affect measurability of the function.

Therefore, it remains to show measurability, for fixed $d \in \mathbb{T}$ and $x' \in X$, of:

$$\{w \in W, s \in S : \phi(d, (x', 0), w, s) \in (x, E)\} \quad (3.3)$$

At the i^{th} internal time (iteration) step, we calculate the increase in elapsed time by scaling w_i by a factor proportional to the total magnitude of the reaction rates. The form of the scaling function, as shown in Algorithm 2.7, is as follows:

$$\tilde{w}_i = \frac{1}{\lambda(x_{t_i})} \times w_i$$

Note that the map from w_i to \tilde{w}_i (defined separately for each i), is therefore continuous over its domain ($\lambda(x_{t_i})$ can only take positive values).

Let us introduce some new notation:

$$G_{k,I} = \{w \in W : \sum_{i=1}^k \tilde{w}_i \in I\}, \quad k \in \mathbb{T}, I \subseteq \mathbb{R}^+ \quad (3.4)$$

$$F_{k,x'} = \{s \in S : \phi(k, x', s, \cdot) = (x, \cdot)\}, \quad k \in \mathbb{T}, x' \in X \quad (3.5)$$

We observe that 3.3 describes the set $G_{d,E} \times F_{k,x'}$. It remains to show each of these sets is measurable.

We prove measurability of $G_{d,E}$ by induction. First assume $d = 1$, and pick an arbitrary interval $I = (a, b)$ (with a or b possibly equal to $\pm\infty$):

$$G_{1,I} = \{w : (\tilde{w}_1 > a \cap \tilde{w}_1 < b)\}$$

Continuity of each transformation $w_i \mapsto \tilde{w}_i$ shows that the set of w_1 for which the above is true is Borel-measurable. In particular, this property holds for $I = E$, as choice of interval was arbitrary.

Now let us assume that $G_{k-1,I}$ is measurable, for some $k \in \mathbb{N}$, and any open interval $I = (a, b)$:

$$G_{k,I} = \bigcup_{q \in \mathbb{Q}} \{G_{k-1,(a-q,\infty)} \cap \{w_k : w_k \in (q, \infty)\}\} \bigcap \bigcup_{q \in \mathbb{Q}} \{G_{k-1,(-\infty,b)} \cap \{w_k : w_k \in (-\infty, q)\}\}$$

We can see that this is a composition of countable unions and intersections applied to measurable sets, and is therefore itself measurable, thus completing our induction argument.

We now prove measurability of $F_{k,x'}$:

Pick $s \in S$. Note that the elements $\{s_i\}_{i=k+1}^\infty$ of s are independent of whether s is in the set $F_{k,x'}$, which is solely a function of the values $\{s_i\}_{i=1}^k$. After all, the substrate vector of a Gillespie simulation at the k^{th} iteration step is independent of the reactions occurring during subsequent iteration steps.

Now observe that, due to the fact that a given element $s_j \in s$ belongs to one of m possible values, there are at most m^k possible distinct sequences $\{s_i\}_{i=1}^k$. Thus there must be a finite number (possibly zero) of distinct specifications on the set $\{s_i\}_{i=1}^k$ that uniquely define any sequence starting with these values to belong to $F_{k,x'}$. Recall from Definition 3.8 the following:

$$B(s, 2^{-k}) = \left\{ q \in S : \{q_i\}_{i=1}^k = \{s_i\}_{i=1}^k \right\}$$

We see that $F_{k,x'}$ is consequently a finite union of open balls in S , which is measurable by definition, with respect to $\mathcal{B}(S)$. □

Corollary 3.13. *Theorem 3.7 is true*

Definition 3.14. *We define the family of shift maps $\theta(s) : \Omega \rightarrow \Omega$ as follows:*

$$\theta(s)(u, w) = (u', w')$$

Here, (u', w') is the pair of sequences satisfying the property:

$$(u'_i, w'_i) = (u_{i+s}, w_{i+s}) \quad \forall i \in \mathbb{N}$$

Theorem 3.15. *The Flow map, φ (Definition 3.5), satisfies the cocycle property over $\theta(\cdot)$ (Definition 3.1).*

Proof. We consider the set of mappings $\varphi(t, w) = \varphi(t, w, \cdot) : Y \rightarrow Y$.

$$\varphi(0, w) = id_Y \tag{3.6}$$

This requires no explanation. A cursory observation of the Gillespie algorithm shows that a 0-step Gillespie simulation is equivalent to the identity mapping. It remains to show the following:

$$\varphi(t+s, \omega) = \varphi(t, \theta(s)\omega) \circ \varphi(s, \omega) \quad \forall s, t \in \mathbb{T} \quad \omega \in \Omega \quad (3.7)$$

The LHS of Equation 3.7 defines a mapping from a system state y_0 to the final value of its $t+s$ step Gillespie simulation, y_{t+s} . The RHS defines the composition of a t step and s step mapping on y_0 . Since a Gillespie simulation proceeds algorithmically, with the results of each iteration stage independent of previous Gillespie states, the RHS also defines a $t+s$ step mapping. Noting Proposition 2.12, we merely need show that the sequences ω_L and ω_R used to perform the simulations on the LHS and RHS of Equation (3.7) respectively, are identical.

From the definition of the flow map, $\omega_L = \omega$, where ω is chosen in Equation (3.7).

On inspection, we see that:

$$\begin{aligned} (\omega_R)_{i=1}^s &= (\omega)_{i=1}^s \\ (\omega_R)_{i=s+1}^\infty &= \theta(s)(\omega)_{i=1}^\infty = (\omega)_{i=s+1}^\infty \end{aligned}$$

So $\omega_L = \omega_R$, and we are done. □

Corollary 3.16. *The flow map φ defines a measurable Random Dynamical System*

Definition 3.17. Generator of the flow We define the generating functions $\zeta(\omega)$ of the flow map as follows:

$$\zeta(\omega) = \varphi(1, \omega) : Y \mapsto Y$$

Current literature focuses on Random Dynamical Systems in which the generating function, as defined above, is an affine map, with respect to the random vector ω that it is parameterised by. This is obviously not the case in our situation, which leaves us at a dead end in terms of our current method of analysis, pending further research on less well-behaved Random Dynamical Systems. This inclusion of this chapter was not just in the hope that such research will be carried out, however. Many of the structures and functions we have defined are subsequently used in other situations during the course of the essay.

4. TIME STABILITY AND ITS EFFECT ON GILLESPIE SIMULATION DYNAMICS

We now introduce a regulatory condition on a Gillespie Model that we name *time stability*, which allows us to define a measure proportional to the limiting probability density of a Gillespie simulation on its substrate state space. First we must attach a probability measure to the space of Gillespie simulations in order that we can quantify simulation properties in terms of their likelihood of occurrence.

Definition 4.1. Simulation State Space

We define the simulation state space of y , O_y , as the space of possible infinitely long Gillespie simulations with initial system state $y \in Y$. We define the generic simulation state space, O , as

$$O = \bigcup_{y \in Y} O_y$$

Recall (Proposition 2.12), that a simulation $o \in O_y$ is determined completely by the element $\omega \in \Omega$. Therefore we use the notation $o_y(\omega)$ to indicate the Gillespie simulation in the space O_y , determined by the random sequence element ω . It is important to recall that the determination is not unique, so it is possible to have distinct elements $\omega_1, \omega_2 \in \Omega$, such that $o_y(\omega_1) = o_y(\omega_2)$.

We define the **k-Step** simulation state space of y , O_y^k , as the space of possible k -step Gillespie simulations with initial system state y .

We will often solely concern ourselves with the substrate vector associated to an element of a simulation $o \in O_y$. The notation $o_i|_X$, signifies the substrate vector associated with o_i , the i^{th} system state of a Gillespie simulation.

Definition 4.2. σ -algebra on O_y

We define a measure, \mathcal{O}_y , on O_y as follows:

O_y^k consists of a set of sequences of system states, of length k . It is therefore isomorphic to the product of k copies of the system state space Y . We have already defined a measure \mathcal{B} on Y (see Definition 3.3). We define \mathcal{O}_y^k to be the product of k copies of \mathcal{B} . We then take:

$$\mathcal{O}_y = \bigcap_{i=1}^{\infty} \mathcal{O}_y^i$$

\mathcal{O}_y , as a product of σ -algebras, is one itself [RUD87].

Definition 4.3. We define the **decomposition map** $C_y : O_y \rightarrow \mathcal{F}$ as follows:

$$C_y(o) = \{\omega \in \Omega : \varphi(i, y, \omega) = o_i \quad \forall i \in \mathbb{T}\}$$

(The construction of the Borel σ -algebra, \mathcal{F} , is found in Definition 3.4)

The decomposition map is named as such due to the fact that it ‘decomposes’ a Gillespie simulation into the set of possible random vector sequences that would have, when inputted into the Gillespie algorithm, provided the original simulation as an output. It can be considered a ‘pseudo’-inverse of the flow map, whose non-injectivity prohibits a true inverse.

Lemma 4.4. *The range of the decomposition map is indeed \mathcal{F} , the Borel σ -algebra on Ω .*

Proof. Pick $o \in O_y$. Let A_i be the set of $\omega \in \Omega$ such that: $\varphi(i, y, \omega) = o_i$. Then we have:

- (1) o_i is measurable in (Y, \mathcal{B})
- (2) Measurability of φ , the flow map implies A_i is measurable

Therefore, we have measurability of $C_y(o) = \bigcap_{i=1}^{\infty} A_i$ □

We proceed to prove measurability of the C_y . We furthermore prove that the image of any \mathcal{O}_y -measurable set, under C_y , is \mathcal{F} -measurable. This implies that the σ -algebra generated by the pre-images of measurable sets in \mathcal{F} , under the decomposition map, is precisely \mathcal{O}_y . Consequently, we can ‘transport’ our existing probability measure \mathcal{V} (see Definition 3.4) on the random sequence space, to a measure on O_y , through this bijective correspondence.

Lemma 4.5. *\mathcal{O}_y -measurability of the decomposition map*

Proof. It is sufficient to prove measurability of the pre-image of any member of a generating set of \mathcal{F} . Recall that any element $\omega \in \Omega$ can be expressed as a pair (u, w) . We can see from 3.4 that \mathcal{F} is generated by sets of the form:

$$\{(u, w) : u_i \in J_1, w_i \in J_2, i \in \mathbb{N}\}$$

where $J_1 \in [0, 1], J_2 \in (0, \infty)$, are open (or closed) intervals.

Pick an arbitrary set $A \in \mathcal{F}$ of this form:

$$A = \{(u, w) : u_k \in J_1, w_k \in J_2, i\}$$

for open intervals J_1, J_2 as previously described.

$$C_y^{-1}(A) = \{o : \exists (u, w) \in A : \varphi(i, y, u, w) = o_i \quad \forall i \in \mathbb{N}\} \tag{4.1}$$

$$= \{o : o_k \in \varphi(k, y, u, w); \quad u_k \in J_1, w_k \in J_2\} \tag{4.2}$$

$$= \{o : (x_k, t_k) \in \varphi(k, y, u, w); \quad u_k \in J_1, w_k \in J_2\} \tag{4.3}$$

Here we recall that $o = \{(x_i, t_i)\}_{i=1}^{\infty}$

It is sufficient to show that the set of (x_k, t_k) satisfying the constraint of Equation 4.3 is \mathcal{B} -measurable, since \mathcal{O}_y is the countable product of copies of \mathcal{B} .

\mathcal{B} is the product space of the σ -algebras generated by the counting measure on X , and the Lebesgue measure on \mathbb{R}^+ . Every subset of X is measurable with respect to the counting

measure, as it is a space of countable cardinality. So we need only prove that the set of $\{t_k\}$ satisfying the constraint of 4.3 is Lebesgue measurable.

Referring to 2.7, we see that:

$$t_k = \tilde{w}_k + \sum_{i=1}^{k-1} \tilde{w}_i$$

where $\tilde{w}_i = \lambda_i w_i$ for some nonzero constant λ_i . Since the support of w_i is $(0, \infty)$, and the set $\{w_i\}_{i=1}^{k-1}$ is unconstrained in Equation 4.3, we can see that the support of $\sum_{i=1}^{k-1} \tilde{w}_i$ is also $(0, \infty)$. Therefore the set of $\{t_k\}$ satisfying the constraints of Equation 4.3 is equivalent to the following set:

$$\{g + h : g \in J_2, h \in (0, \infty)\}$$

This is an interval, so Lebesgue measurable. \square

Lemma 4.6. *The image of an \mathcal{O}_y -measurable set under C_y is \mathcal{F} -measurable*

Proof. \mathcal{O}_y is the product of countable copies of \mathcal{B} , the σ -algebra on Y , the system state space. Therefore, any \mathcal{O} -measurable set A will be of the form:

$$A = \prod_{i=1}^{\infty} A_i : \quad A_i \in \mathcal{B} \quad \forall i$$

Now by definition:

$$C_y(A) = \bigcap_{i=1}^{\infty} \{\omega : \varphi(i, y, \omega) \in A_i\}$$

Each element of the intersection is \mathcal{F} -measurable due to measurability of the flow map φ (Theorem 3.7). Therefore, $C_y(A)$ is \mathcal{F} -measurable. \square

Definition 4.7. Probability measure on \mathcal{O}_y

We define the function $\mathcal{P} : \mathcal{O}_y \mapsto [0, 1]$ as follows:

$$\mathcal{P}(o) = \mathcal{V}(C_y(o))$$

The bijection under C_y of members of \mathcal{O}_y and \mathcal{F} , (Lemmas 4.5, 4.6), combined with the fact the \mathcal{V} is a probability measure, assures us that \mathcal{P} is a probability measure.

We can now analyse the dynamics of Gillespie simulations probabilistically. In a biochemical system of the type the Gillespie algorithm attempts to model, we inevitably observe stability of the dynamics over time, despite their local stochasticity. We proceed to define this idea rigorously, and demonstrate its consequences.

Definition 4.8. Time-Stable simulation state space

For a given $x \in X$, $y \in Y$, let

$$A_{xy} = \left\{ o \in O_y : \bigcap_{n=1}^{\infty} \left(\bigcup_{m=n}^{\infty} \{o_m : o_m|_X = x\} \right) \right\} \quad (4.4)$$

In other words, A_{xy} is the set of simulations in O_y that return to the substrate vector x infinitely often. We can see that Equation 4.4 is measurable. We define O_y to be time-stable (with respect to x) if there exists $x \in X$ such that

$$\mathcal{P}(A_{xy}) = 1$$

Lemma 4.9. Suppose O_y is time-stable with respect to some substrate vector $x' \in X$. Then O_y is time stable with respect to any $x \in X$. We may therefore alter our definition of time stability to remove dependence on a given substrate vector $x \in X$.

Proof. For any $o \in A_{x'y}$, we have a subsequence $\{o_{n_i}\}_{i=1}^{\infty}$ of elements of o whose substrate vector is equal to x' . Let K_i be the event that there exists $j \in \{n_i, \dots, n_{i+1}\}$ such that $o_j|_X = x$. We can see that $\{K_i\}$ are independent, identically distributed random variables, due to the following:

- (1) The evolution of the substrate vector in the Gillespie algorithm is dependent only on its current state (ie it is a Markov chain), due to independence of of sequence elements of sequences in the random sequence space.
- (2) The only system state K_i is dependent on is therefore o_{n_i} , whose substrate vector is x , for all i

The same reasoning assures us that the set of σ -algebras: $\{\mathcal{O}_y^{n_i} \cap (\mathcal{O}_y^{n_{i-1}})^c\}_{i=1}^{\infty}$ are mutually independent. We also see that

$$K_i \in \mathcal{O}_y^{n_i} \cap (\mathcal{O}_y^{n_{i-1}})^c$$

Let

$$G_k = \sigma \left(\bigcup_{k=i}^{\infty} \{\mathcal{O}_y^{n_i} \cap (\mathcal{O}_y^{n_{i-1}})^c\} \right)$$

Now let K be the event that, for any $i \in \mathbb{N}$, with $i > n_1$, there exists $j \in \mathbb{N}$ such that $j > i$ and $o_j = x$, with $j \geq n_1$. Then

$$K = A_{xy} = \bigcap_{i=1}^{\infty} \left(\bigcup_{j=n_i}^{\infty} K_j \right) \in \bigcap_{i=1}^{\infty} G_i$$

So A_{xy} is a tail event with respect to the aforementioned set of σ -algebras, and, by Kolmogorov's Zero-One Law [WIL91], has probability 0 or 1.

Now given $o \in A_{x'y}$, $i \in \mathbb{N}$, we know that there exists j such that $n_j > i$. In addition, $o_{n_j}|_X = x'$ by definition. Due to the assumptions in Definition 2.3, we are assured that there is a finite sequence of reactions that would take the substrate vector from x' to x .

The probability of each of these reactions occurring is always strictly greater than zero, regardless of the system state, again by assumption. Therefore the product of these positive probabilities, which gives the probability of the aforementioned sequence of reactions occurring, is again greater than 0. This event would imply the occurrence of an element of o with substrate vector x' , and index greater than n_j . This is an event contained within K , and therefore, A_{xy} . This shows that $P(A_{xy}) > 0$, so $P(A_{xy}) = 1$. \square

The previous lemma shows that, given a time stable simulation state space, any Gillespie simulation will hit every substrate vector infinitely often with probability one. One would think that this does not make sense in the context of a biological system. However, note that the above proof gives no constraints on the time scale required to achieve such time stability.

Corollary 4.10. *Suppose there exists a time stable simulation state space, O_y on a Gillespie model (X, R) . Then every simulation state space on (X, R) is time stable, regardless of initial system state. This allows us to generalise the notion of time stability from a property of a particular simulation state space to one of the whole Gillespie model.*

Proof. Suppose O_y is time stable. Given arbitrary $y' = (x', t') \in Y$, we are required to show that $O_{y'}$ is time stable. Now we know from Lemma 4.9 that, on O_y , $\mathcal{P}(A_{x'y}) = 1$. Let

$$H = \{o \in O_y : \exists k \text{ s.t. } o_k|_X = x'\}$$

Time stability ensures that $\mathcal{P}(H) = 1$. Therefore $\mathcal{P}(H \cap A_{x'y}) = 1$. So any simulation in O_y that hits x' , returns to hit x' again infinitely often, with probability one. But this is equivalent to saying that any simulation with initial substrate vector x' , returns to hit x' again infinitely often, with probability one, due to memorylessness of the Gillespie algorithm. This implies that $O_{y'}$ is time stable. \square

We proceed to prove the almost -sure existence of a limiting probability for elements of a Gillespie simulation on a time stable model to take a given substrate state. This then allows us to define a measure on the substrate state space proportional to the aforementioned limiting probability.

Definition 4.11. *Given a time stable simulation space O_y , an arbitrary fixed $x \in X$, and an element $o \in A_{xy}$ containing a subsequence $\{o_{n_i}\}_{i=1}^{\infty}$ with $o_{n_i}|_X = x$, we define*

$$N_i^o = n_i^o - n_{i-1}^o; \quad N_1^o = n_1^o$$

*Note that these random variable N_i^o are independent and identically distributed (with respect to both o and i), by a similar argument to that of the independence of $\{K_i\}$ in Lemma 4.9. So their distribution is solely a function of x , and not O_y . We take the random variable N^x to be a representative random variable of the set. In other words it shares their distribution. We term the distribution of N^x to be the **return distribution** of x .*

Lemma 4.12. *For a given time stable Gillespie model, if the mean of $N^{x'}$, $\mu_{x'}$, is finite for a given $x' \in X$, then the mean of N^x is finite for every $x \in X$. We refer to such a model as **strongly time stable**. Any time stable Gillespie model that does not satisfy this property is **weakly time stable***

Proof. Let $C_{x'x}$ be the random variable denoting the number of steps a Gillespie simulation starting at x' takes to reach x . Define $C_{xx'}$ similarly. Then we see that

$$\mathbb{E}(N^x) \leq \mathbb{E}(C_{x'x}) + \mathbb{E}(C_{xx'}) + \mathbb{E}(C_{x'x})$$

(1) $\mathbb{E}(C_{xx'}) < \infty$:

For a simulation o with subsequence $\{o_{n_i}\}$ of system states with substrate vector x' , let A be the event that, for a given i , there exists $k \in [n_i, n_{i+1}]$ such that $o_k|_X = x$, with p its corresponding probability. Then

$$p \times \mathbb{E}(N^{x'}|A) \leq \mu_{x'}$$

The reasoning is thus: Let us denote the probability space of $N^{x'}$ as (Ω, \mathcal{F}, P) . Then

$$\mu_{x'} = \int_{\Omega} N^{x'} dP = \int_A N^{x'} dP + \int_{\Omega/A} N^{x'} dP$$

All terms are positive due to non-negativity of the support of $N^{x'}$. Now if $\int_A N^{x'} dP \geq \frac{1}{p} \times \mu_{x'}$, we have:

$$\int_A N^{x'} dP \geq p \times \int_A \mu_{x'} dP = p \times \frac{1}{p} \times \mu_{x'} \quad \perp$$

Now $C_{xx'} \leq \{N^{x'}|A\}$ by definition, so the result follows.

(2) $\mathbb{E}(C_{x'x}) < \infty$:

$$\mathbb{E}(C_{x'x}) = \sum_{i=1}^{\infty} i \times p_i$$

Here, p_i is the probability that a Gillespie simulation with initial substrate vector x' has a substrate vector x after i steps. Let k be the probability that the simulation returns to x' without hitting x . Then $k < 1$, by time stability of the model. We also have

$$\sum_{i > n \times \mu_{x'}} p_i < k^n$$

Now take $\tilde{p} = \sup_{i < \mu_{x'}} (p_i)$. Then

$$\mathbb{E}(C_{x'x}) \leq \left(\sum_{i=1}^{\infty} k^i \right) \left(\sum_{i \leq \mu_{x'}} i \times \tilde{p} \right) < \infty$$

□

Theorem 4.13. *Let O_y be a simulation state space on a time stable Gillespie Model (X, R) . Fix an arbitrary $x \in X$, and, for a given $o \in O_y$, let*

$$B_n^{o,x} = \left\{ \frac{1}{n} \times \#\{0 \leq k < n : o_k|_X = x\} \right\}$$

Suppose that the model is strongly time stable, so that N^x has finite mean, μ . Recall that μ is independent of the initial state y of O_y . Then we claim that

$$\mathcal{P} \left(\left\{ o : \liminf_{n \rightarrow \infty} B_n^{o,x} = \limsup_{n \rightarrow \infty} B_n^{o,x} = \frac{1}{\mu} \right\} \right) = 1$$

In other words, the subset of O_y in which $\lim_{n \rightarrow \infty} B_n^o$ exists and equals $\frac{1}{\mu}$ has probability one.

Now suppose N^x has infinite mean. Then

$$\mathcal{P} (\{ \liminf_{n \rightarrow \infty} B_n^{o,x} = 0 \}) = 1$$

Proof.

We deal with the case where N^x has finite mean first:

We are assured that A_{xy} is a set of measure one, by definition. For every $o \in A_{xy}$, there exists $\{o_{n_i}\}_{i=1}^{\infty}$ with $o_{n_i}|_X = x$. Therefore, on this set, by the strong law of large numbers:

$$\mathcal{P} \left(A_{xy} \cap \left\{ o : \lim_{i \rightarrow \infty} \frac{1}{i} \sum_{k=1}^i N_k^x = \mu \right\}^C \right) = 0$$

This implies

$$\mathcal{P} \left(\left\{ o : \lim_{i \rightarrow \infty} \frac{1}{i} \sum_{k=1}^i N_k^x = \mu \right\} \right) = 1$$

Now

$$\left\{ \lim_{j \rightarrow \infty} B_j^{o,x} = \frac{1}{\mu} \right\} = \left\{ \lim_{j \rightarrow \infty} \frac{1}{j} n_j^o = \mu \right\} = \left\{ \lim_{j \rightarrow \infty} \frac{1}{j} \sum_{i=1}^j N_i^o = \mu \right\}$$

Thus proving the theorem.

The proof is more involved where N^x has infinite mean:

A sufficient condition for the claim is that

$$\forall \epsilon > 0 : \mathcal{P} \left(\left\{ o : \bigcap_{i=1}^{\infty} \bigcup_{j=i}^{\infty} \{B_j^o > \epsilon\} \right\} \right) = 0 \tag{4.5}$$

Now

$$\{B_j^o > \epsilon\} = \{n_{c_j}^o < \epsilon j\} = \left\{ \sum_{i=1}^{c_j} N_i^o < \epsilon j \right\}$$

where $c_j = \lceil \epsilon j \rceil + 1$, the value of ϵj rounded up to the nearest positive integer. Therefore (4.5) is equivalent to

$$\forall \epsilon > 0 : \mathcal{P} \left(\left\{ o : \lim_{i \rightarrow \infty} \bigcup_{j=i}^{\infty} \left\{ \sum_{k=1}^{c_j} N_k^o < \epsilon j \right\} \right\} \right) = 0 \quad (4.6)$$

We note that $\epsilon j \leq c_j$. Therefore (4.6) is implied by:

$$\forall \epsilon > 0 : \mathcal{P} \left(\left\{ o : \lim_{i \rightarrow \infty} \bigcup_{j=i}^{\infty} \left\{ \frac{1}{c_j} \sum_{k=1}^{c_j} N_k^o < 1 \right\} \right\} \right) = 0 \quad (4.7)$$

Regardless of the value of c_j , the random variable $\tilde{N}_{j,k}^o = \frac{1}{c_j} N_k^o$ will have infinite mean, since this property is invariant under scaling. With our new notation, (4.7) is equivalent to

$$\forall \epsilon > 0 : \mathcal{P} \left(\left\{ o : \lim_{i \rightarrow \infty} \bigcup_{j=i}^{\infty} \left\{ \sum_{k=1}^{c_j} \tilde{N}_{k,j}^o < 1 \right\} \right\} \right) = 0 \quad (4.8)$$

A sufficient condition for the above is that there exists some value of j with $\tilde{N}_{k,j}^o > 1$, with probability one. Such an occurrence is a tail event, by the same reasoning as given in Lemma 4.9 regarding the set K . Therefore it has either probability 0 or 1. Since any element $\tilde{N}_{k,j}^o > 1$ has infinite mean, there is a strictly positive probability that it will be greater than one, ruling out the former option. We are done. \square

Lemma 4.14. *A Random variable with non-negative support must have defined mean, whether finite or infinite.*

Proof. The expected value of a random variable X on a probability space $(\Omega, \mathcal{F}, \mathcal{P})$ is defined as the Lebesgue integral:

$$\int_{\Omega} X d\mathcal{P}$$

This does not exist only when both the positive and negative components of the integral both sum to infinity. However, if X has non-negative support, the above integral has no negative component. \square

Corollary 4.15. *Any substrate vector on a time-stable Gillespie model must satisfy one of the conditions of Theorem 4.13. So any time stable Gillespie model is either strongly or weakly time stable.*

Definition 4.16. Given a strongly time stable Gillespie Model, we attach a measure $\mathcal{L} : X \rightarrow [0, 1]$ as follows [FAY95]:

$$\mathcal{L}(x) = \frac{1}{\mu_x}$$

Here, μ_x is the mean value of N^x , the return distribution of x . Theorem 4.13 and Corollary 4.15 show us that it is defined for all $x \in X$. Since it is just a weighted form of the counting measure, we neglect to prove that it does indeed define a measure. We name \mathcal{L} the **Limiting Probability** of X . Note that this is a probability measure. By Theorem 4.13:

$$\mathcal{P} \left(\bigcap_{x \in X} \{o : \liminf_{n \rightarrow \infty} B_n^{o,x} = \limsup_{n \rightarrow \infty} B_n^{o,x}\} \right) = 1$$

Taking a simulation o from the above set, we have:

$$\mathcal{L}(X) = \mathcal{P} \left(\left\{ \lim_{n \rightarrow \infty} \bigcup_{x \in X} B_n^{o,x} \right\} \right) = 1$$

5. A TIME STABILITY HEURISTIC

Notation 5.1.

(1) We take $d : X \times X \mapsto \mathbb{R}^+$ to be the L_1 metric on X . So

$$d(x, x') = \sum_{i=1}^n |x_i - x'_i|$$

(2) We take $B(x, r)$ to define a closed ball with respect to the L_1 metric. So

$$B(x, r) = \{x' \in X : d(x, x') \leq r\}$$

Definition 5.2. Rests

Consider a Gillespie model (X, R) , with $A \subset X$. Suppose we have a simulation $o \in O$, with a finite string of elements: $\{o_i\}_{i=k}^m$.

- (1) $o_k|_X \notin A$
- (2) $o_i|_X \in A \forall i = \{k+1, \dots, m\}$
- (3) $o_{m+1} \notin A$

Then we define the string $\{o_i\}_{i=k+1}^m$ to be a **rest** of o in A , of length $m - k$. Suppose we replace conditions (2) and (3) with:

- (4) $o_i|_X \in A \forall i > k$

Then we refer to the string $\{o_i\}_{i=k+1}^\infty$ as an **infinite rest** of o in A .

Define $x \in A$ as the substrate vector: $o_k|_X$. We refer to x as the **entry point** of the rest.

Lemma 5.3. *Suppose we have two results of a subset A , with identical entry point $x \in A$. Then the random variables denoting resting length are independent and identically distributed, with defined (possibly infinite) mean.*

Proof. Independence and identical distribution are a consequence of the memorylessness of the Gillespie Algorithm. The defined mean is a consequence of 4.14 \square

Definition 5.4. Stable Mass Model

We define a Gillespie Model (X, R) to be of stable mass if there exists a subset $N \subset X$ of the form $N = B(x, r)$ such that the following condition is satisfied for all Gillespie simulations o :

$$\exists c \in (0.5, 1) : \text{ s.t. } o_k|_X \notin N \Rightarrow \mathcal{P}\left(d(o_{k+1}|_X, N) < d(o_k|_X, N)\right) > c$$

Here, we assume that the reactions of the system satisfy the assumptions of Definition 2.3.

Theorem 5.5. *A Gillespie Model (X, R) of stable mass is also strongly time stable*

Proof. We know that the set N is of the form $B(x, r)$ by definition. Let $y = (x, 0)$. For a simulation $o \in O_y$, let us denote by Q_o the expected resting length of o in X/N . We prove this is finite in the next lemma. Assuming for now that it is indeed finite, we know that there exists $x \in N$ such that the expected resting time of o in X/x is finite, due to the fact that N contains a finite number of substrate vectors. But this is equivalent to stating that the expected value of the return distribution of x is finite, which implies that the Gillespie model is strongly time stable. \square

Lemma 5.6. *The expected resting time of a simulation o in X/N exists and is finite in a stable mass model, where $o_1|_X \in N$.*

Proof. We know that the set N is of the form $B(x, r)$ by definition. Let $y = (x, 0)$. Then, for a simulation $o \in O_y$, let us define the function:

$$f : O_y \mapsto \mathbb{Z}^N; \quad f(o)_k = d(o_k|_X, x)$$

Here, \mathbb{Z}^N is the sequence space on \mathbb{Z} . We can see that f is measurable. Also;

$$f(o)_k > r \Leftrightarrow o_k|_X \notin N$$

Given a rest of o in X/N : $\{o_i\}_{i=j}^{j+m}$, with unknown length m , we wish to find the expected length m of the rest, assuming that it exists. This transforms to the problem: Given the information that $f(o)_k = r + 1$, what is the expected first return time of $\{f(o)_i\}$ to the

interval $[0, r]$, assuming that it exists?

Define the random variables $\{Z_i\}_{i=k}^m$ as follows:

$$Z_i = \begin{cases} 1 & f(o)_i \geq f(o)_{i=1} \\ -1 & f(o)_i < f(o)_{i=1} \\ 1 & i = k \end{cases}$$

Then $S_j = \sum_{i=k}^j Z_i$ is a biased random walk starting at 1. The time of the first expected hitting of S_j to the point 0 is always greater or equal to the expected value of m , as S_j is by definition greater or equal to $f(o)_j - f(o)_k$. As i tends to infinity, S_i will, with probability one, eventually be smaller or equal to \tilde{S}_i , where:

$$\tilde{S}_j = \sum_{i=k}^j \tilde{Z}_i; \quad \tilde{Z}_i = \begin{cases} 1 & \text{with probability } 1-c \\ -1 & \text{with probability } c \\ 1 & \text{if } i = k \end{cases}$$

This is due to the fact that the probability of the event $\{Z_i = -1\}$ is always greater than c , by the definition of mass stability. We now show that a first hitting time exists with probability one for $\{\tilde{S}_j\}$, which implies that it exists for $\{S_j\}$ with probability one as well: Calculation gives us the following:

$$\mathbb{E}(\tilde{S}_n) = 1 + (n-1)(1-2c)$$

$$Var(\tilde{S}_n) = 1 + (n-1)(1-c)(1+c)4c$$

Applying Chebyshev's Inequality [WIL91], gives:

$$P\left(|\tilde{S}_n - \mathbb{E}(\tilde{S}_n)| \geq k \times \sqrt{Var(\tilde{S}_n)}\right) < \frac{1}{k^2}$$

Since $\mathbb{E}(S_n)$ decreases at a rate proportional to n (since $1-2c < 0$), while $\sqrt{Var(S_n)}$ increases at a rate proportional to the square root of n , one can see that $\{\tilde{S}_n\}$ will almost surely hit 0 eventually, as n increases. So m exists with probability one.

It remains to show that the expected value of this first hitting time, for the sequence $\{\tilde{S}_i\}$, is finite. This implies that the same is true for $\{S_i\}$, which in turn shows that the expected value of m is finite.

Let H be the random variable denoting the length of steps required for the random walk \tilde{S}_i to hit zero, and p_i be the probability of the event $\{H = i\}$ (which is only nonzero when i is odd). Then

$$p_{2i+1} = \binom{2i+1}{i} c^{i+1} (1-c)^i$$

$$\mathbb{E}(H) = \sum_{i=0}^{\infty} i p_i \approx \sum_{i=1}^{\infty} (2i+1) \left(\frac{2i+1}{2i}\right)^{\frac{3}{2}} c (4c(1-c))^i \frac{1}{\sqrt{\pi n}} \quad (5.1)$$

Equation 5.1 uses Stirling's approximation. We note that convergence of the approximation is sufficient to prove that of the original summation, due to the sum of error terms being bounded. Now $4c(1-c) < 1$, due to the fact that $c \in (0.5, 1)$. The presence of this exponentially decreasing term assures us that the sum converges. \square

Example 5.7. Comparison of Time Stability of Random Walks

Pólya proved that a random walk on a one or two dimensional integer lattice will return to its starting point with probability one, but that in higher dimensions, the probability of return was strictly smaller than one. Also proven was that the expected return time of the one and two dimensional random walks to their starting point was infinite [POL21]. Subsequent research estimated the probability of return on a three dimensional integer lattice to be approximately 0.34 [GLA77]. In Appendix A, we have formulated Gillespie Models equivalent to the one and three dimensional random walks, just described. With respect to our terminology, the first is weakly time stable, while the second is not time stable at all.

We construct and prove strong time stability of a weakly constrained three dimensional random walk Gillespie model, in Appendix A.4, using Theorem 5.5. The counterintuitive physical interpretation is that an unbiased random walk will meet a given point on the \mathbb{Z}^3 lattice with probability less or equal to 0.35, while one in which a weak drift term towards the origin is present outside of a ball around the origin, will meet every point with probability one.

Example 5.8. A Circadian Model

Consider our Circadian Model described in Appendix B, and graphed in Figure 3. The probability of a reaction of decreasing mass, given substrate vector $x \in X$ is equal to:

$$\sum_{i=1}^{10} \frac{R_{i+10}(x)}{R_i(x) + R_{i+10}(x)}$$

A sufficient condition for mass stability of the model around a set of the form $B(0, r)$, for some finite r , is that the above summation is always greater than some $c > 0.5$, when x is outside of the set $B(0, r)$. Upon algebraic manipulation, this reduces to the following inequality being satisfied, for some $K_r > 0$, dependent on r :

$$(v_2 - v_3)x_1 + v_4x_{10} - \frac{v_1}{1 + x_{10}^4} > K \sum_{i=1}^{10} x_i$$

This inequality is not satisfied mathematically on all but a set of finite mass, as would be required for mass stability around a set of the form previously described. It fails due to the fact that the ratio between the quantities $\sum_{i=2}^9 x_i$ and x_1 , is unbounded. However, incorporation of a bound in the Gillespie model, which would be physically realistic, given that x_1 levels drive the rate of production of $\{x_2, \dots, x_9\}$, would imply mass stability of the model.

Time Stability with respect to External Time

Time stability is defined with respect to internal time; the number of system state changes a Gillespie undergoes. One easily observes that it is preserved under the transformation to external time. What is not so clear, however, is whether strong time stability is invariant under this transformation. In other words, does the (in)finiteness of the expected value of the internal time return distribution determine that of its external counterpart?

Conjecture 5.9. *There exist Gillespie models that are only weakly time stable, such that the expected external time between a simulation hitting the same substrate vector is finite.*

Reasoning Consider the Modified Random walk given in Appendix A.2. The dynamics are the same as in the unmodified version with respect to internal timestep, and it is therefore weakly time stable. However, the amount of elapsed time between reactions decreases doubly exponentially as the mass of the system increases. It could therefore be possible that the external expected return time between oscillations was finite. A sufficient condition for this to be true would be that there existed an interval $I = [-r, r]$, for some $r \in \mathbb{N}$, such that the expected elapsed resting time of \mathbb{Z}/I was finite. A heuristic for the proof would be to show the terms in the expansion of the expected value of this resting time decreased fast enough as r increased, to allow the above to be true for some finite value of r . This would then prove it to be true for any value of r .

Note that computer simulation may not be a useful tool in the determination of time stability. A seemingly time stable system might in fact have a critical substrate state, unlikely to be reached after a computationally tractable length of time, which, when reached, has a probability of irreversibly changing the dynamics of a simulation in such a way that previously attained substrate states are no longer reached with probability one.

6. A CATEGORISATION OF THE CONCEPT OF OSCILLATORY DYNAMICS ON A GILLESPIE MODEL

The notion of an oscillating system is well defined in the deterministic setting. One merely need find a system state that is repeated after some defined time value, which we call the period of the oscillation. Adding an external noise to a deterministic oscillation does not complicate the definition unduly. The underlying trajectory would be unaffected by the noise values in such a setting, and could easily be uncovered by spectral analysis. In a Gillespie simulation, however, the path of the underlying trajectory will be determined by past noise values, which makes determination of period, and, indeed, what constitutes an oscillatory behaviour, ill defined. In this section we outline some possible means of arriving to such a definition.

Proposition 6.1. *Any strongly time stable Gillespie model oscillates in the sense that there is an expected (finite) return time of a simulation on the model to any substrate vector it has already reached.*

For a proof of the above proposition, refer to Section 4. Strong time stability, however, is not a sufficient condition for the existence of what we would empirically term an oscillatory model. For instance, a random walk on a finite lattice will revisit every point infinitely often, with a finite expected return time. We would however hardly characterise this as an oscillatory system.

The above proposition shows that the notion of eventual return to the same system state is ill-suited to the stochastic setting. Let us therefore consider another property of deterministic oscillators. A sine wave could be considered the most basic oscillatory system. The

function:

$$f(x) = \sin(x)$$

has a period of 2π . We could describe any point on this trajectory by its phase, which would generally be defined as its value modulus 2π . An alternative definition of phase could be given by constructing a bijective mapping between the phase as previously described, and some other property of the trajectory. We do so by fixing a timestep α such that the quantity $f(t + \alpha) - f(t)$ is unique to the value of t , modulus 2π . One appropriate value of α is one that corresponds to one quarter of the period. Figure 1 graphs $f(t + \alpha)$ against $f(t)$. We see that this results in a perfect circle. We can now define the phase of a point $f(t)$ by the angle of the point it corresponds to on the circle.

In a similar vein, a Gillespie simulation of the Circadian Model (Appendix B), part of which is plotted in Figure 2, was realised computationally. An algorithm (omitted) was created to determine an average period of oscillation, β , for the PER1 protein. We use the word oscillation loosely here, as, since we have not yet constructed an exact definition. The algorithm determined the average spacing between peaks, where multiple vacillations in quick succession around a single local maximum were ignored. Let $g(t)$ be the trajectory of the Gillespie simulation. Figure 4 provides a graph of $g(t + \beta)$ against $g(t)$. We note that the shape of the graph bears only a very superficial resemblance to the circle obtained under the analogous transformation of the sine function (Figure 1). However, a shared property of the graphs is the constant unidirectional change in angle with respect to a central point. This corresponds to the ‘phase’ of each model increasing in the same direction with time. We now construct a possible definition of a stochastic oscillator that is based on the existence of this property.

Definition 6.2. Phase Function

Pick $c \in \mathbb{R}^n$. We define the phase function of c as follows:

$$Q_c^n : \mathbb{R}^n \times \mathbb{R}^n \mapsto [0, 2\pi)^n$$

$$(Q_c^n)_i(a, b) = \text{atan2}((a_i - c_i), (b_i - c_i)) \quad i \in \{1, \dots, n\}$$

Here, *atan2* is a commonly used modification of the arctangent function. *atan2(a, b)* gives the angle, in radians, between the positive *x*-axis and the vector from the origin to the Cartesian co-ordinate (a, b) . [MAT12].

Definition 6.3. Jump Function

Consider a Gillespie Model (X, R) together with a positive integer α . Pick an $x \in X$, and take $y = (x, 0) \in Y$. Given $o \in O_y$, we take:

$$J_{\alpha, o} : X \mapsto X$$

$$J_{\alpha, o}(x) = o_\alpha|_X$$

Definition 6.4. Mapping to the Phase Space

Consider a Gillespie Model (X, R) , where $X \cong \mathbb{N}^n$. Pick a point $c \in \mathbb{R}^n/X$, a positive

integer $\alpha \in \mathbb{N}$, and a Gillespie simulation o with $o_1|_X = x$. We define a mapping $T_{c,\alpha}$ from $X \times O_{(x,0)}$ to $[0, 2\pi)^n$ as follows:

$$T_{c,\alpha}(x, o) = Q_c^n(x, J_{\alpha,o}(x))$$

We also define:

$$H_{c,\alpha}(x) = \mathbb{E}(T(x)) = \int_{o \in O_{(x,0)}} T_{c,\alpha}(x, o) \, d\mathcal{P}$$

Our expected value is taken with respect to the probability measure espoused in Definition 4.7. We refer to $\text{Im}(H_{c,\alpha})$ as the **phase space** with respect to the parameters of H .

The i^{th} component of $T_{c,\alpha}(x, o)$, to put it plainly, calculates the difference in angle between the two lines originating at c_i , and respectively ending at x_i and the position of x_i after α steps of the simulation o . $\mathbb{E}(T(x_i))$ gives the expected value of this angle over the probability space of simulations. We require that both lines are non-trivial (ie not a point), in order for such an angle to exist. The constraint of c to \mathbb{R}^n/X ensures this.

H , meanwhile, gives a family of mappings from the substrate vector space X to the phase space $[0, 2\pi)^n$, parameterised by the values c and α . Note that, since there are only a countable number of elements in X , any H -mapping is necessarily non-surjective, and any phase space is countable.

Definition 6.5. Generators on the Phase Space

Recall the flow map φ , from Definition 3.5. We define a family of mappings on the phase space, parameterised by values $c \in \mathbb{R}^n/X$, and $\alpha \in \mathbb{N}$, as follows:

$$K_{c,\alpha} : \text{Im}(H_{c,\alpha}) \mapsto [0, 2\pi)^n$$

$$K_{c,\alpha}(h) = \left(\mathbb{E} \left(H_{c,\alpha}(\varphi|_X(1, \omega, y)) | H_{c,\alpha}(y|_X) = h \right) - h \right) \quad \text{mod } 2\pi$$

Note that the expected value here is taken with respect to the probability measure \mathcal{V} defined on Ω in Definition 3.4.

Given that a system state is in a given phase, with respect to some c and α , the function $K_{c,\alpha}$ outputs the expected value of the change in phase after one step of the Gillespie Algorithm. In an oscillatory system, we would want the direction of this change in phase to be invariant with respect to the original phase. In other words, we want K to be an order preserving map. We give a rigorous definition of an order preserving map below:

Definition 6.6. Order Preserving Map

Suppose we pick three arbitrary vectors $h^1, h^2, h^3 \in \mathcal{K} \subset [0, 2\pi)^n$, with the i^{th} component of the vector h^j denoted h_i^j . We define

$$f : \mathcal{K} \mapsto [0, 2\pi)^n$$

to be an order preserving map if, for all $i \in \{1, \dots, n\}$:

$$h_i^{j_1} > h_i^{j_2} > h_i^{j_3} \Rightarrow$$

- (1) $f(h_i^{j_1}) > f(h_i^{j_2}) > f(h_i^{j_3})$ OR
- (2) $f(h_i^{j_3}) > f(h_i^{j_1}) > f(h_i^{j_2})$ OR
- (3) $f(h_i^{j_2}) > f(h_i^{j_3}) > f(h_i^{j_1})$

In other words, the order of the points on the canonical circle formed by identifying the ends of the interval $[0, 2\pi)$, is preserved.

Definition 6.7. Absolutely Oscillatory System

We define a strongly time-stable Gillespie Model (X, R) to be an absolutely oscillatory system if there exists $c \in \mathbb{R}^n / X$, and $\alpha \in \mathbb{N}$ such that $K_{c,\alpha}$ is an order preserving map. We name $K_{c,\alpha}$ a **rotation map** of the system. Note that it may not be unique.

Note that one could also define a relatively oscillating system by requiring $K_{c,\alpha}$ to be order preserving on a restriction of its domain to a set of high limiting probability.

The obvious way to form a notion of the period of an absolutely oscillatory system would be to find the average ‘rotation’ realised by the rotation map over its domain. This cannot be done by iterating the rotation map infinitely, and finding the average rotation, a technique often used in a dynamical systems setting (See ‘Poincare Rotation Number’, [DEV92]). The reason is that the domain and range of the map are not equal. We instead proceed as follows:

Definition 6.8. Measure on the Phase Space

Assume we have a strongly time stable Gillespie Model. Recall the Limiting probability measure on X defined in 4.16. We can transport it to any phase space of the form $Im(H_{c,\alpha})$ as follows:

$$\mathcal{L}(h) = \mathcal{L}(\{x \in X : H_{c,\alpha}(x) = h\})$$

Countability of both spaces makes measurability of the mapping trivial

Lemma 6.9. Suppose (X, R) is an absolutely oscillatory Gillespie model, with a rotation map $K_{c,\alpha}$. Let us turn the (countable) elements of the phase space into a sequence $\{h_i\}_{i=1}^\infty$. Let us define

$$S_n = \sum_{i=1}^n K_{c,\alpha}(h_i) \mathcal{L}(h_i)$$

Then $\lim_{n \rightarrow \infty} S_n$ exists and is invariant with respect to permutations of the indexing of the set $\{h_i\}$

Proof.

$$\sum_m^\infty S_m \leq 2\pi \sum_m^\infty \mathcal{L}(h_m) \leq 2\pi$$

Therefore $\{S_n\}$ is a Cauchy sequence and convergent. Since every term in the summation is non-negative, it is also absolutely convergent, and the limit is therefore independent of the indexing of $\{h_i\}$. \square

Definition 6.10. Period of the Rotation Map of an Absolutely Oscillatory System

Given a rotation map $K_{c,\alpha}$, and taking S_n as in Lemma 6.9, we define the period of the rotation map, $\beta_{c,\alpha}$ of an oscillator as

$$\lim_{n \rightarrow \infty} S_n$$

Lemma 6.11.

Suppose $K_{c,\alpha}$ is a rotation map. For a simulation o , let

$$\tilde{S}_k^o = \frac{1}{k} \sum_{i=2}^k (H_{c,\alpha}(o_i) - H_{c,\alpha}(o_{i-1}))$$

Then

$$\mathcal{P}(\{o : \lim_{k \rightarrow \infty} \tilde{S}_k^o = \beta_{c,\alpha}\}) = 1$$

Proof. Given a subsequence $\{o_{n_i}\}$ such that $H_{c,\alpha}(o_{n_i}|_X) = h_j$, we are assured, by the strong law of large numbers, that

$$\lim_{k \rightarrow \infty} \frac{1}{k} \sum_{n_i=2}^k (H_{c,\alpha}(o_{n_i}) - H_{c,\alpha}(o_{n_{i-1}})) = K_{c,\alpha}(h_j)$$

with probability one. The result is then a consequence of Theorem 4.13. \square

Conjecture 6.12. Any two rotation maps of an absolutely oscillatory Gillespie Model have the same period

Reasoning

Given two rotation maps $K_{c,\alpha}$ and $K_{c',\alpha'}$, we would expect, by the strong law of large numbers, that ;

$$\mathcal{P}\left(\left\{o \in O_y : \left(\lim_{k \rightarrow \infty} \frac{1}{k} \sum_{i=1}^k K_{c,\alpha}(o_i|_X) = \beta_{c,\alpha}\right) \cap \left(\lim_{k \rightarrow \infty} \frac{1}{k} \sum_{i=1}^k K_{c',\alpha'}(o_i|_X) = \beta_{c',\alpha'}\right)\right\}\right) = 1$$

for any initial system state $y \in Y$. If $\beta_{c,\alpha}$ and $\beta_{c',\alpha'}$ were different, then the above set of simulations would ‘wrap around’ the points c and c' at different rates, in the respective phase spaces. This would be equivalent to the limiting rate of oscillation of a simulation being different with respect to two different viewing points on the substrate state space, which would only be possible if both viewing points were not ‘centred’ inside the expected oscillation trajectory. However, the placing of the points c and c' are both necessarily ‘centred’ in such a way, in order for $K_{c,\alpha}$ and $K_{c',\alpha'}$ to define rotation maps.

APPENDIX A. RANDOM WALK MODELS

A.1. One Dimensional Random Walk.

The following Gillespie Model (X, R) recreates a random walk on the integer lattice:

- (1) $X \cong \mathbb{Z}$
- (2) Reactions r_1, r_2 satisfy assumptions of 2.3
- (3) $R_1(x) = R_2(x) = 0.5$

A.2. Modified One Dimensional Random Walk.

The following Gillespie Model (X, R) recreates a random walk on the integer lattice:

- (1) $X \cong \mathbb{Z}$
- (2) Reactions r_1, r_2 satisfy assumptions of 2.3
- (3) $R_1(x) = R_2(x) = 0.5(\exp(\exp(\sum_{i=1}^n x_i)))$

A.3. Three Dimensional Random Walk.

The following Gillespie Model (X, R) recreates a random walk on the three dimensional integer lattice:

- (1) $X \cong \mathbb{Z}^3$
- (2) Reactions r_1, r_2, \dots, r_6 satisfy assumptions of 2.3
- (3) $R_i(x) = \frac{1}{6} \quad \forall i \in \{1, 2, \dots, 6\}$

A.4. Three Dimensional Weakly Constrained Random Walk. The following Gillespie Model (X, R) recreates a weakly constrained random walk on the three dimensional integer lattice:

- (1) $X \cong \mathbb{Z}^3$
- (2) Reactions r_1, r_2, \dots, r_6 satisfy assumptions of 2.3
- (3) For some $r \in \mathbb{N}, d \in (0.5, 1)$, we have:

$$R_i(x) = \begin{cases} \frac{1}{6} & x \in B(0, r) \\ \frac{1-d}{3} & x \notin B(0, r), i \in \{1, 2, 3\} \\ \frac{d}{3} & x \notin B(0, r), i \in \{4, 5, 6\} \end{cases}$$

Lemma A.1. *The Three Dimensional Weakly Constrained Random Walk Model is strongly time stable*

Proof. By Theorem 5.5, it is sufficient to prove that the model is of stable mass. But this is obvious upon observation of the definition of stable mass; for any Gillespie simulation $o \in O$, we have:

$$o_k|_X \notin N \Rightarrow \mathcal{P}(d(o_{k+1}|_X, N) < d(o_k|_X, N)) = 3 \left(\frac{d}{3} \right) = d > c$$

for some $c > 0.5$.

□

APPENDIX B. THE CIRCADIAN MODEL

The existence of the circadian clock, nearly universal to animal species, is hypothesised to derive from a complex autoregulatory gene network, which undergoes one complete oscillation over an approximately 24 hour period [GOLO3]. The expression of the protein PER1 are understood to be involved in this process. Goldbeter et al. [GOL03] propose a 30 parameter model of the dynamics of PER1 expression levels, which they then realise through the Gillespie Algorithm. We simplified this model, while aiming to reflect the overall dynamics of the process. A summary of the simplified processes affecting PER1 regulation is described below, and derived from [MOR08] (in which an alternative simplified model is suggested):

- (1) The presence of PER1 mRNA in the cell ribosome instigates translation of a protein we shall term IP1 ('Intermediate Protein 1'). This is essentially a completely de-phosphorylated version of the PER1 protein.
- (2) IP1 undergoes several stages of phosphorylation, at a constant rate that we shall term the *protein forward rate*, into IP2, IP3, etc.
- (3) The final product realised after complete phosphorylation is the PER1 protein. The presence of PER1, which acts as a transcription factor, inhibits transcription of the PER1 mRNA.

Note that both the PER1 protein and PER1 mRNA in this process degrade constantly at a rate proportional to their overall level in the cell.

Upon computational experimentation, I determined that at least 8 intermediate phosphorylation stages were required to induce an obviously oscillatory Gillespie Algorithm realisation of the model. The Gillespie Model (X, R) used is described below, and a graphical simulation is provided (Figure 3):

- CONSTANTS: *(Experimentally derived in [GOLO3])*
- (1) Ω (Scaling constant, setting size of the system)
 - (2) $v_1 = 30\Omega$ (MRNA transcription rate)
 - (3) $v_2 = 36$ (MRNA degradation rate)
 - (4) $v_3 = 2$ (MRNA translation rate)
 - (5) $v_4 = 0.5$ (Transcription Factor degradation rate)
 - (6) $\lambda = 1.666667$ (Phosphorylation rate)

SUBSTRATES:

X_1 describes the quantity of PER1 mRNA in the ribosome.

X_2, \dots, X_9 describe the quantities of the intermediate proteins in the nucleus.

X_{10} describes the quantity of the PER1 protein in the nucleus.

$$X = (X_1, X_2, \dots, X_{10}) \cong \mathbb{N}^{10}$$

REACTIONS

The reactions of the system, $\{r_1, r_2, \dots, r_{20}\}$, follow the assumptions of Definition 2.3. The reaction rate function, R , is as follows:

$$\begin{aligned} R_1(x) &= \frac{v_1}{1 + x_{10}^4} & R_{11}(x) &= v_2 \times x_1 \\ R_2(x) &= v_3 \times x_1 & R_i(x) &= \lambda \times x_{i-10} \quad i \in \{12, \dots, 19\} \\ R_i(x) &= \lambda \times x_{i-1} \quad i \in \{3, 4, \dots, 10\} & R_{20}(x) &= v_4 \times x_{10} \end{aligned}$$

APPENDIX C. DETERMINISTIC ANALOGUES

Definition C.1. *Deterministic Analogue*

Consider a Gillespie Model (X, R) , with n substrate and m reaction types. Recall that $X \cong \mathbb{N}^n$.

We define the deterministic analogue of the Gillespie Model (X, R) to be a pair $(\mathbb{R}^n, \tilde{R})$, satisfying the properties:

- (1) $\tilde{R} : \mathbb{R}^n \rightarrow \mathbb{R}^m$
- (2) $\tilde{R}(x) = R(x) \quad \forall x \in \mathbb{R}^n \cap X$
- (3) Given any initial value problem of the form:

$$x'(t) = \tilde{R}(x(t)) \quad (x_0, t_0) \in (\mathbb{R}^n \cap X, \mathbb{R}^+)$$

$$\text{We have: } x_i(t) \geq 0 \quad \forall t > 0, i \in 1 \dots n$$

We say \tilde{R} is an extension of R to \mathbb{R}^n . It is not unique.

We define the deterministic analogue of a Gillespie trajectory, with initial substrate vector x_0 and elapsed time t_0 , to be a solution of the Initial Value Problem:

$$x'(t) = \tilde{R}(x(t))$$

Definition C.1 gives the sense that choosing an appropriate deterministic analogue to a Gillespie Model presents a problem. In practice, most Gillespie Models are derived from pre-existing deterministic analogues, rather than vice versa, sidestepping the problem. In any case, a computationally practicable Gillespie model is likely to have a closed-form expression for the reaction rate vector field, which will admit an obvious continuous extension, as can be seen in, for example, the Circadian Model described in Appendix B.

Conjecture C.2. *Given any Gillespie Model (X, R) , there exists a deterministic analogue $(\mathbb{R}^n, \tilde{R})$ such that \tilde{R} is both continuous and infinitely differentiable.*

A proof of the conjecture is outside the scope of the essay. Suffice to say that, given n points in \mathbb{R}^n , it is possible to construct a degree n polynomial passing through each of them. Similarly, given a countable, nowhere-dense, subset of \mathbb{R}^n , it is possible to construct a power series function, which is by definition continuous and infinitely differentiable, that passes through each point.

APPENDIX D. INDEX OF TERMINOLOGY

Substrate State Space:	<i>Definition 2.1</i>
System State Space:	<i>Definition 2.2</i>
Reaction State Space:	<i>Definition 2.3</i>
Reaction Rate Vector Field:	<i>Definition 2.4</i>
Gillespie Model:	<i>Definition 2.5</i>
Classical Gillespie Algorithm:	<i>Algorithm 2.7</i>
Gillespie Simulation:	<i>Definition 2.8</i>
Trajectory of Gillespie Simulation:	<i>Definition 2.9</i>
Random Sequence Space:	<i>Definition 2.11</i>
Modified Gillespie Algorithm:	<i>Algorithm 2.13</i>
Random Dynamical System:	<i>Definition 3.1</i>
Flow Map:	<i>Definition 3.5</i>
Simulation State Space:	<i>Definition 4.1</i>
Decomposition Map:	<i>Definition 4.3</i>
Time Stability:	<i>Definition 4.8</i>
Return Distribution:	<i>Definition 4.11</i>
Time Stability: Strong, Weak:	<i>Lemma 4.12</i>
Limiting Probability:	<i>Definition 4.16</i>
Rests: Finite, Infinite:	<i>Definition 5.2</i>
Stable Mass Model:	<i>Definition 5.4</i>
Phase Function:	<i>Definition 6.2</i>
Jump Function:	<i>Definition 6.3</i>
Phase Space:	<i>Definition 6.4</i>
Order Preserving Map:	<i>Definition 6.6</i>
Absolutely Oscillatory System, Rotation Map:	<i>Definition 6.7</i>
Period of Rotation Map:	<i>Definition 6.10</i>

APPENDIX E. FIGURES

FIGURE 1. Plot of $\sin(x)$ against $\sin(x + \alpha)$; $\alpha = 0.5\pi$

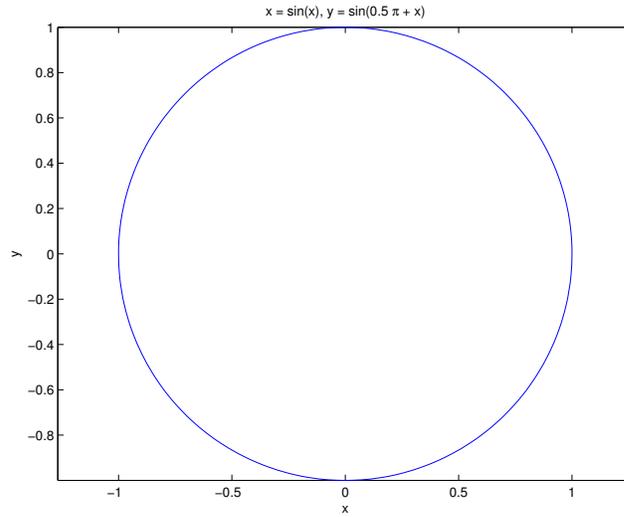


FIGURE 2. A trajectory of the Deterministic Analogue to the Circadian Model described in Appendix B

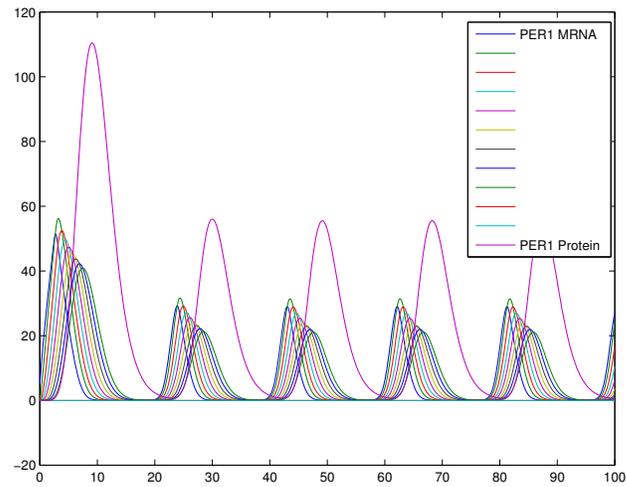


FIGURE 3. A Gillespie trajectory of the Circadian Model described in Appendix B , with scale parameter $\Omega = 1$

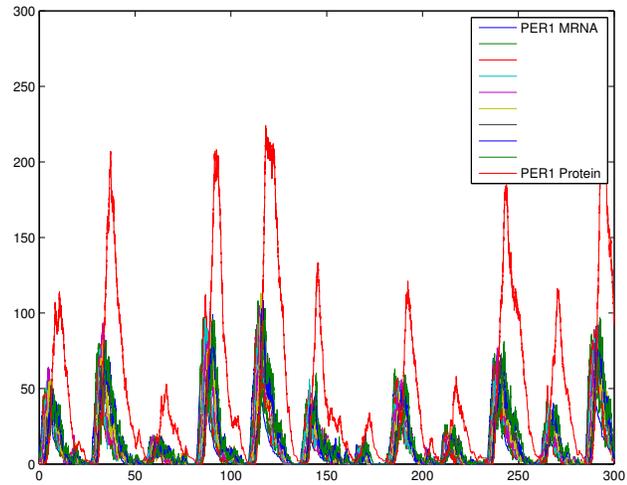
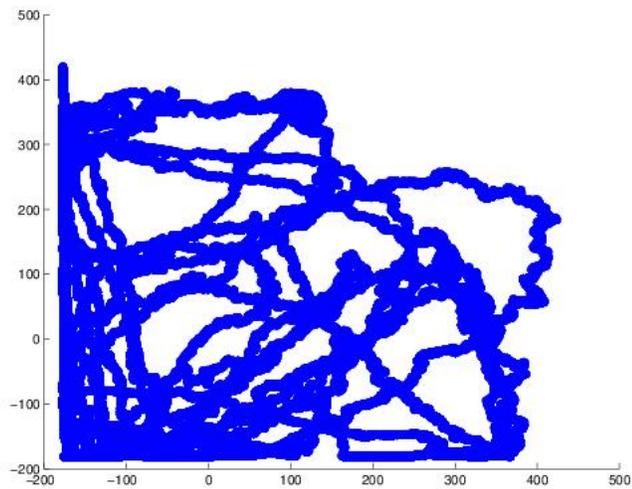


FIGURE 4. Plot of $g(t+\beta)$ against $g(t)$. Here, $g(t)$ is the Gillespie trajectory of PER1 protein in a Gillespie simulation of the Circadian model. $\beta = 9.8$



APPENDIX F. A LIST OF DEFINITIONS AND RESULTS THAT HAVE ANALOGUES IN
MARKOV CHAIN THEORY

Strongly and weakly time stable Gillespie Models correspond to positive recurrent and recurrent ergodic Markov Chains [FAY95]. The construction of these objects involves **Lemma 4.9** and **Corollary 4.10**, which have analogues in [FAY95].

The formulation of **Lemma 4.12** was motivated by a similar result in [FAY95].

Theorem 4.13 corresponds to the Ergodic Theorem for Markov Chains [FAY95]

The idea of **Definition 4.16** was motivated by an identical result in [FAY95]

All of the proofs of the above results were arrived at completely independently of [FAY95], and do not bear a resemblance to their existing counterparts

7. BIBLIOGRAPHY

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